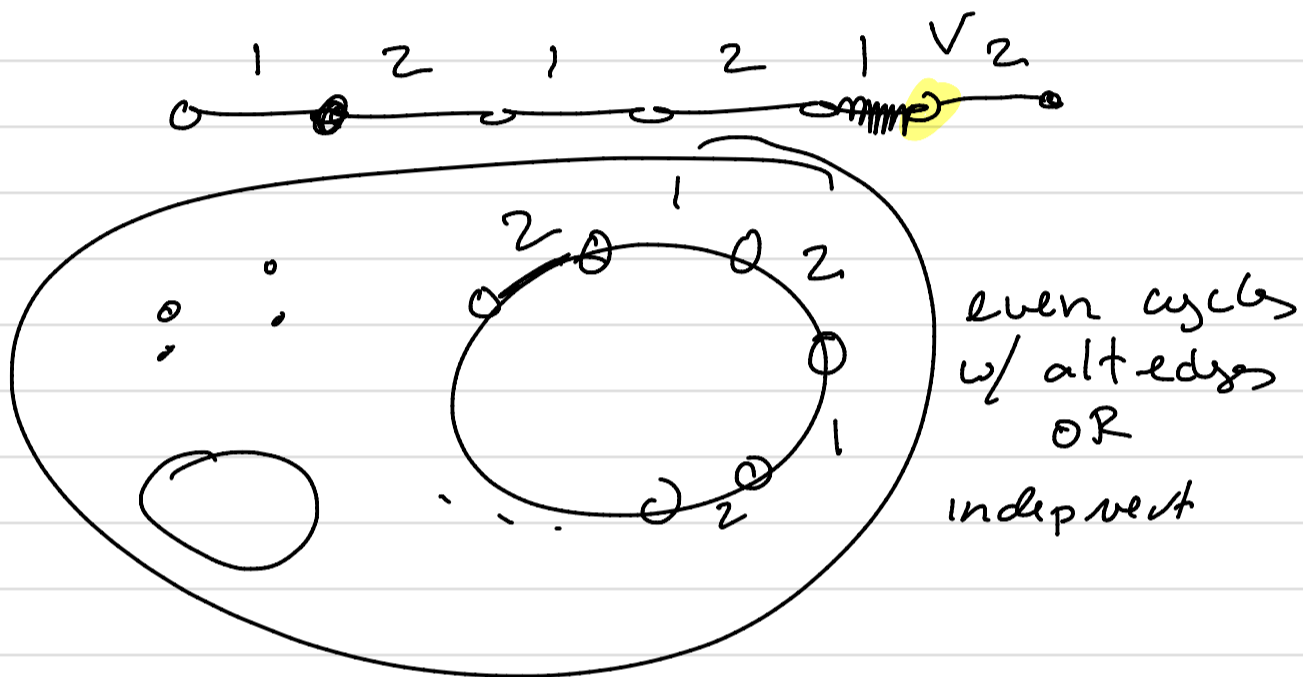


Wed 27 Sept

- Hmwk 4 due Fri
- Hmwk 3 + Soln returned

If  $G$  has two 1-factors,  $M_1$  and  $M_2$ ,  
what does  $M_1 \Delta M_2$  look like?

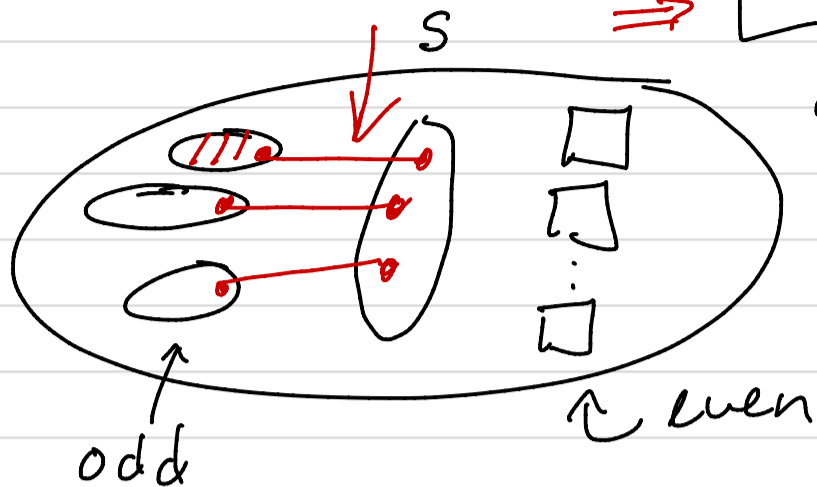


Thm 2.2.1 Tutte's Thm

$G$  has a 1-factor  $\iff$

$\forall S \subseteq V(G), \quad q(G-S) \leq |S|$

# of odd comp's of  $G-S$



$G$

$M$

$S = \emptyset$   
 $q(G-S) = q(G) = 0$

Tutte's cond.  
T's cond

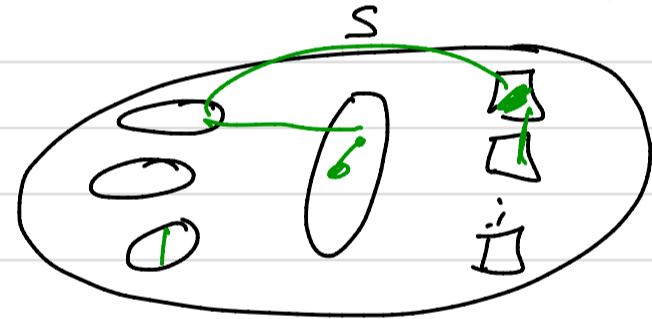
Pf:  $\Leftarrow$ : (by contradiction)

\* Suppose  $\exists G$  s.t.  $G$  satisfies Tutte's cond. and  $G$  has no 1-factor.

• obs:  $|G|$  is even.  $S = \emptyset$

• Choose  $G$  to be edge maximal.

- If  $e \in \bar{G} \wedge G+e$  no 1-factor, then  $G := G+e$
- Stop when no such  $e$  exists, new  $G'$
- $G'$  still satisfies T's cond b/c



$G$

$G'$

adding edges cannot increase  $q(G-S)$

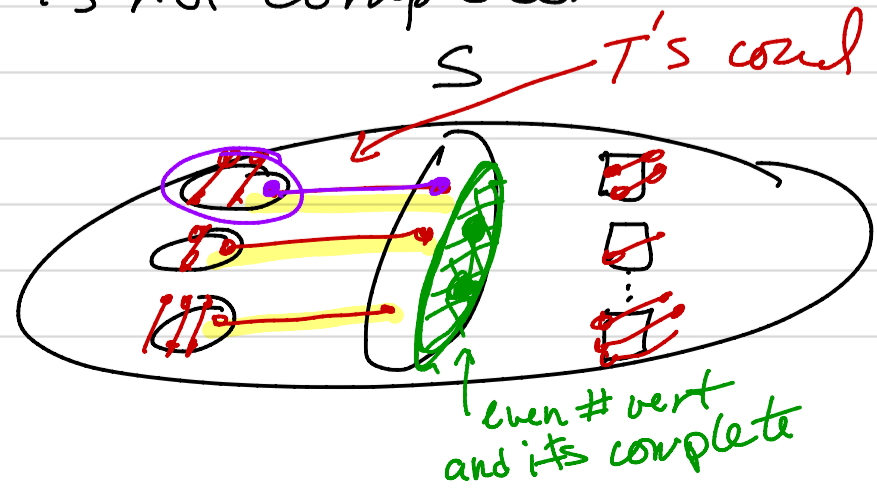
- If  $e \in \bar{G}'$ , then  $G'+e$  has a 1-factor.

• Strategy: Show  $G'$  has a 1-factor  $\implies \Leftarrow$ .

• Choose  $S = \{v \in G' : d_{G'}(v) = |G'|-1\}$

- Case 1: The components of  $G'-S$  are complete
- Case 2: There is 1 comp,  $C$ , of  $G'-S$  that is not complete.

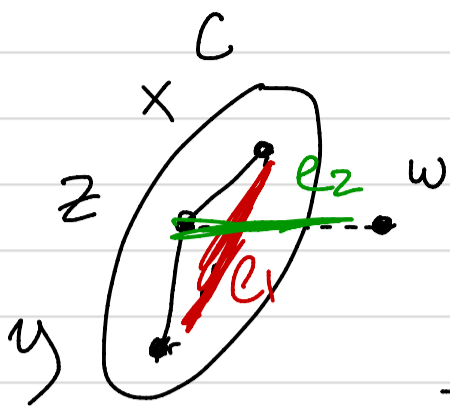
Case 1:



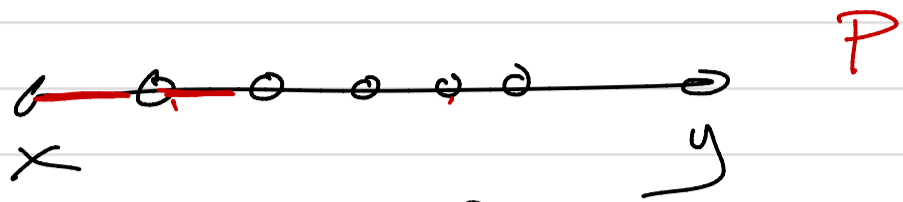
$G$

Find 1-factor  
 $M$

Case 2  $\exists$  a noncomplete conpt  $C$  of  $G'-S$



So  $\exists$  a shortest  $xy$  path in  $C$ ,  $P$   
Then  $P$  must contain an induce  $P^3$



$\exists$  some  $w \in V(G')$  s.t.  $zw \notin E(G')$   
b/c  $z \notin S$ .

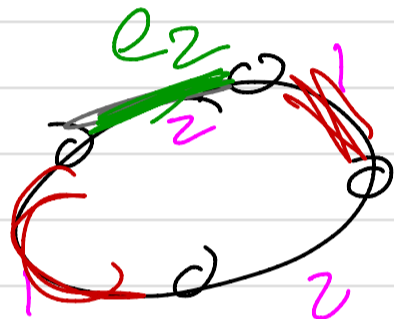
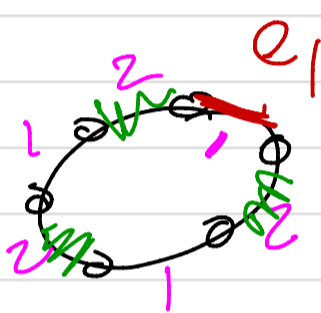
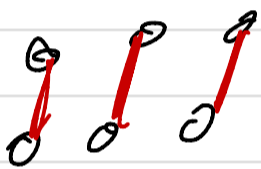
$xy = e_1 \in \overline{G'}$   
 $zw = e_2 \in \overline{G'}$

$G' + e_1$  has 1-factor  $M_1$   
 $G' + e_2$  " " "  $M_2$

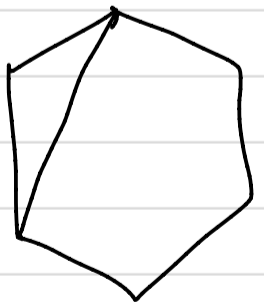
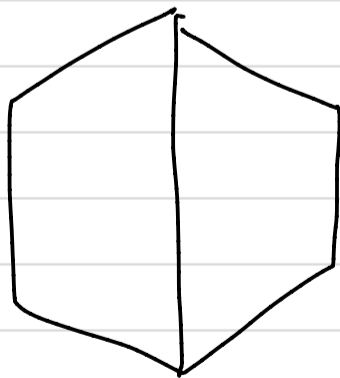
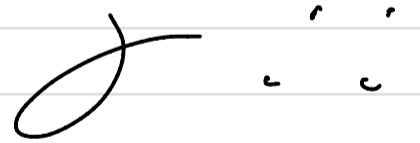
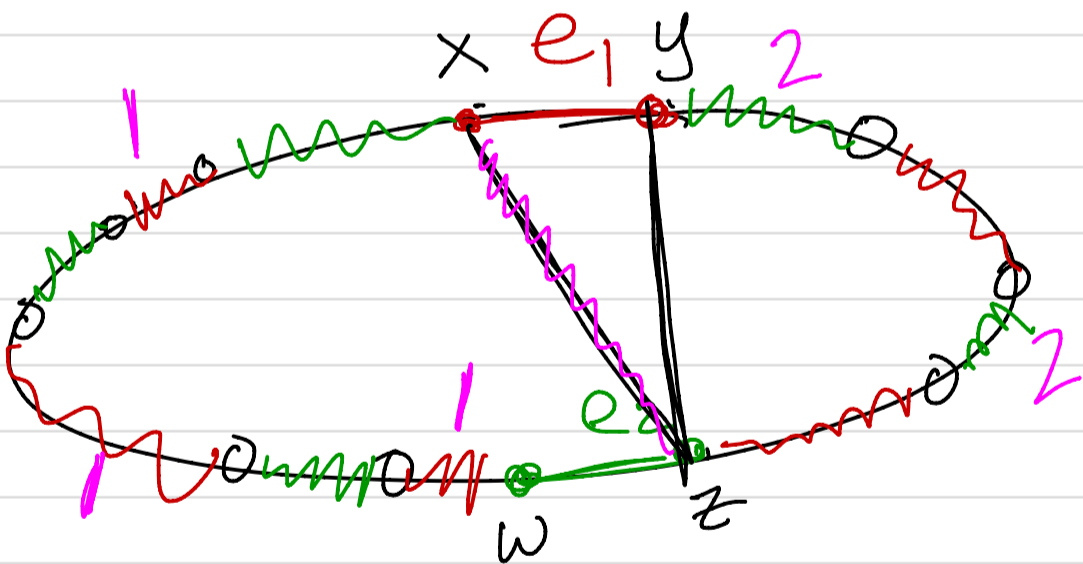
$M_1 \Delta M_2$  is nonempty, b/c  $e_1$  and  $e_2$  are here.

Subcase A

$G'$



Subcase B

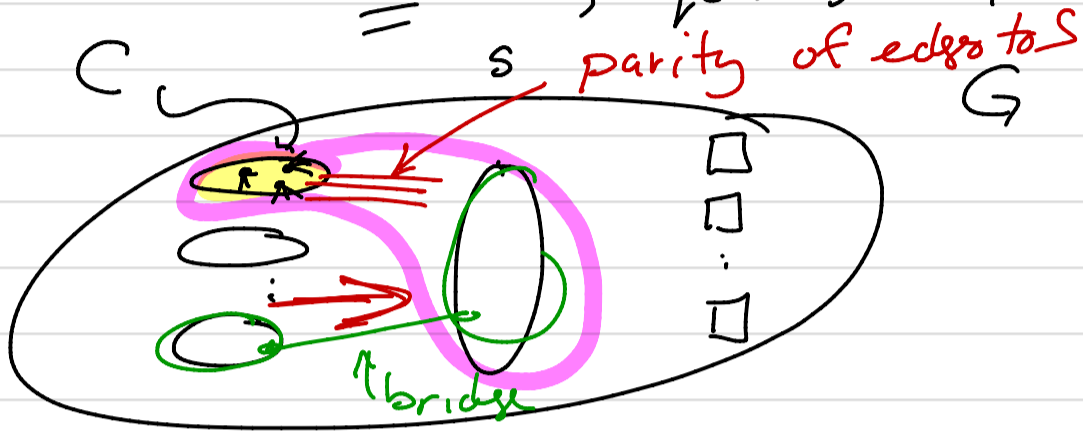


Cor 2.2.2

A bridgeless, cubic graph has a 1-factor.

translation: cubic  $\equiv$  3-regular  
 bridgeless  $\equiv \forall e \in G, G-e$   
 still connected.

Pf: Show  $G$  cubic  $\wedge$  bridgeless  $\Rightarrow$   
 $\forall S \subseteq V, q(G-S) \leq |S|$



$C$  some odd cycle of  $G-S$

$$\sum_{v \in C} d_G(v) = 3|C| \leftarrow \text{odd}$$

$$\sum_{v \in C} d_C(v) \leftarrow \text{even}$$

$$\text{all edges} = \underbrace{\text{edges in } C}_{\text{even}} + \underbrace{\text{edges to } S}_{\text{odd}}$$

$\forall$  odd cycle in  $G-S$ , it sends 3 edges to  $S$ .

$$\# \text{ edges from odd cycles to } S \geq 3 \cdot q(G-S)$$

But  $S$  can accept at most  $3|S|$

$$\text{Now } 3|S| \geq 3 \cdot q(G-S)$$

$$|S| \geq q(G-S)$$

