

Fri 3 Nov

- Homework due today
- Agenda
 - Prove Ford-Fulkerson
 - Start Ch7 Extremal Theory
- Two Weeks to Mid 2.
Should we schedule the Fairbanks
Midterm on Thurs from 2:30-4:30?
- Monday before Thanksgiving will be a Zoom
Lecture.

Thm 6.2.2 Ford-Fulkerson

Given $N = (G, s, t, c)$ network. Let f represent a flow on N .

$$\max_{\substack{f \text{ flow on} \\ N}} |f| \stackrel{\leq}{=} \min_{\substack{S \subseteq V \\ S \text{ a cut in } N}} \{c(S, \bar{S})\}$$

Pf: • It is sufficient to construct f and find S so that $|f| = c(S, \bar{S})$.

- Our construction implies f is integral.
- Start with $f_0(\vec{e}) := 0$. (integral, valid flow)
- Suppose we have flows f_0, f_1, \dots, f_k so that

$$|f_i| + 1 \leq |f_{i+1}|$$

- How to construct f_{k+1} ?

- def: $W = x_0 \vec{e}_0 x_1 \vec{e}_1 x_2 \dots x_{l-1} \vec{e}_{l-1} x_l$ a good walk ↖ sv
if $x_0 = s, x_l = t$ and $\forall 1 \leq i \leq l-1 \quad c(\vec{e}_i) > f(\vec{e}_i)$

$$\forall i, \vec{e}_i = (e_i, x_i, x_{i+1})$$

define

$$S_k = \{s\} \cup \{v \in V : \exists \text{ a good } s v \text{ walk}\}$$

Case 1: $t \in S_k$ (Case 2 $t \notin S_k$)

Pf: \exists a good $s t$ path in N .

$$P: x_0 \vec{e}_0 x_1 \vec{e}_1 \dots x_{l-1} \vec{e}_{l-1} x_l = t$$

Let $\varepsilon = \min_{\vec{e}_i \in \vec{P}} \{c(\vec{e}_i) - f(\vec{e}_i)\} > 0$ and an integer.

\uparrow \uparrow
 \neq \neq

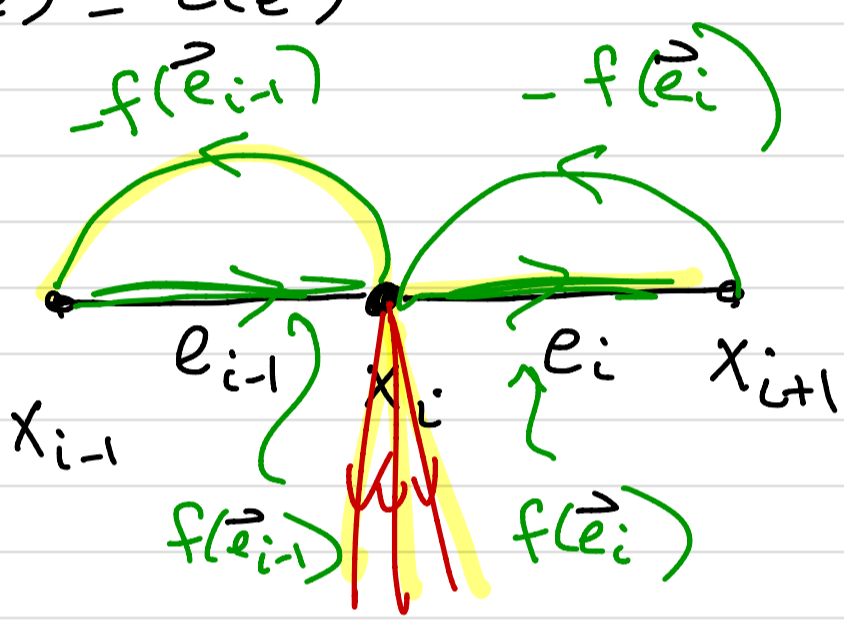
$$f_{k+1}(\vec{e}) = \begin{cases} f_k(\vec{e}) + \epsilon & \text{if } \vec{e} = \vec{e}_i \\ f_k(\vec{e}) - \epsilon & \text{if } \vec{e} = \overleftarrow{e}_i \\ f_k(\vec{e}) & \text{if } e \in E(P) \end{cases}$$

① $f(\vec{e}) = -f(\overleftarrow{e})$

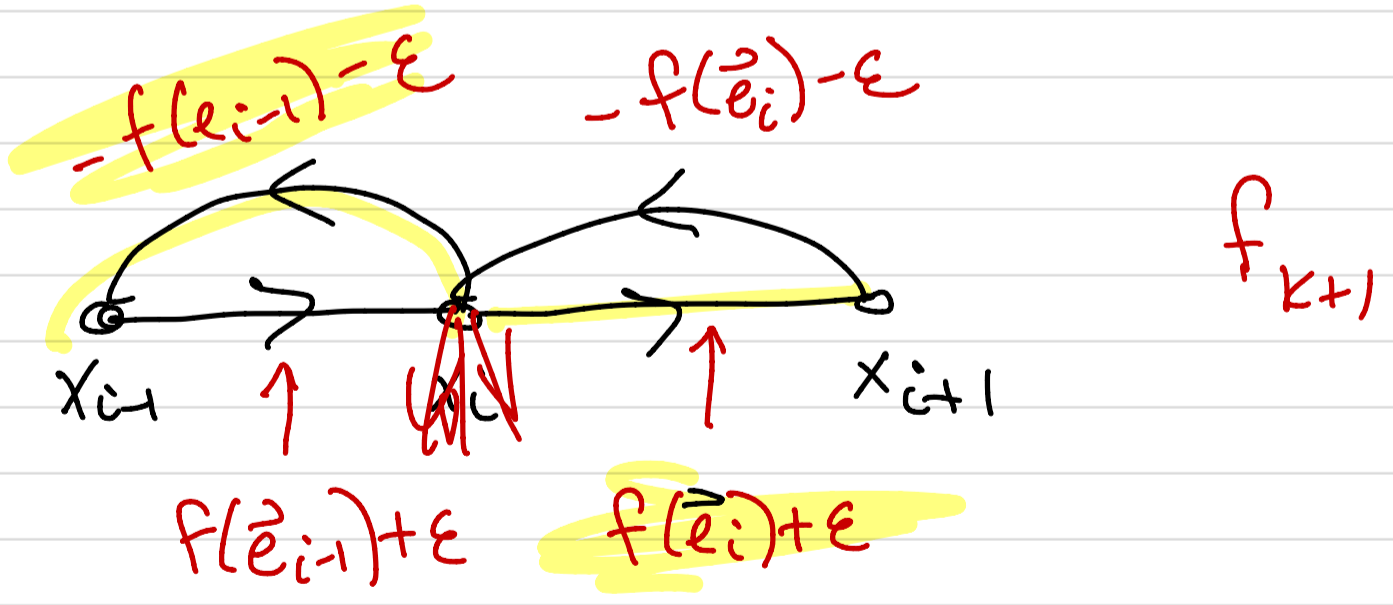
and $|f_{k+1}| \geq |f_k| + 1$

② $f(u, v) = 0$ for $v \in V - \{s, t\}$

③ $\forall \vec{e} \quad f(\vec{e}) \leq c(\vec{e})$



$f(x_i, V) = \text{sum of flow on red edges} + f(e_i) - f(\overleftarrow{e}_{i-1}) = 0$

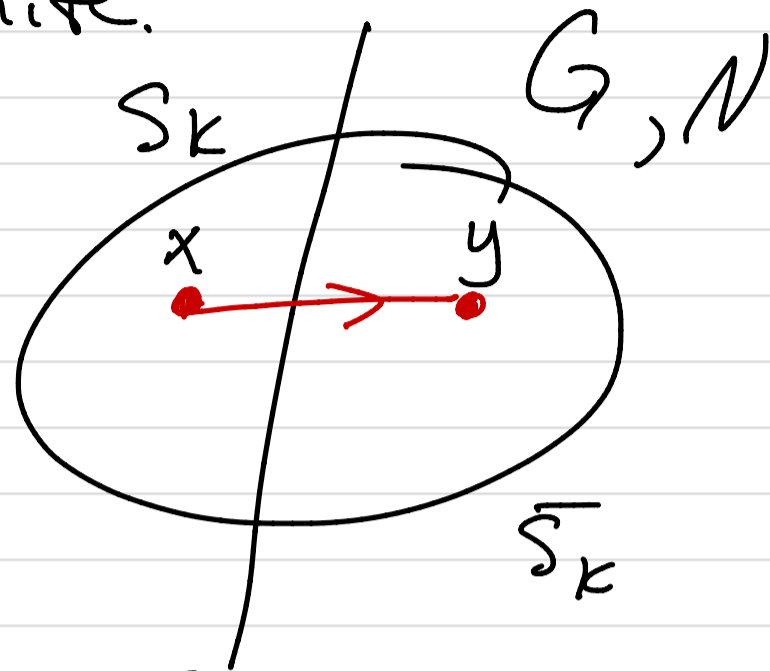


$f_{k+1}(x_i, V) = \text{sum on red edges} + f_k(\vec{e}_i) + \epsilon - f_k(\overleftarrow{e}_{i-1}) - \epsilon = 0$

Obs : This incrementing process must terminate b/c capacities are finite.

Case : $t \notin S_k$

$$|f| = c(S_k, \bar{S}_k)$$



• S_k is a cut b/c $s \in S_k, t \in \bar{S}_k$

• $\forall \vec{e} \in \underline{E(S_k, \bar{S}_k)}, c(\vec{e}) = f(\vec{e})$

$$e = xy, \vec{e} = (e, x, y)$$

otherwise $y \in S_k$.

b/c



Ch 7 Extremal Graph Theory

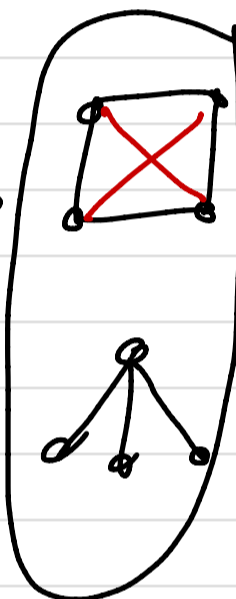
a) Assume # vertices in G is fixed, n .
 How many edges can G have and
 still not have a K^3 ?

Obs) Such a graph is necessarily
 K^3 -critical, (ie $\forall e \in \bar{G}$,
 $G+e$ has a K^3 .)

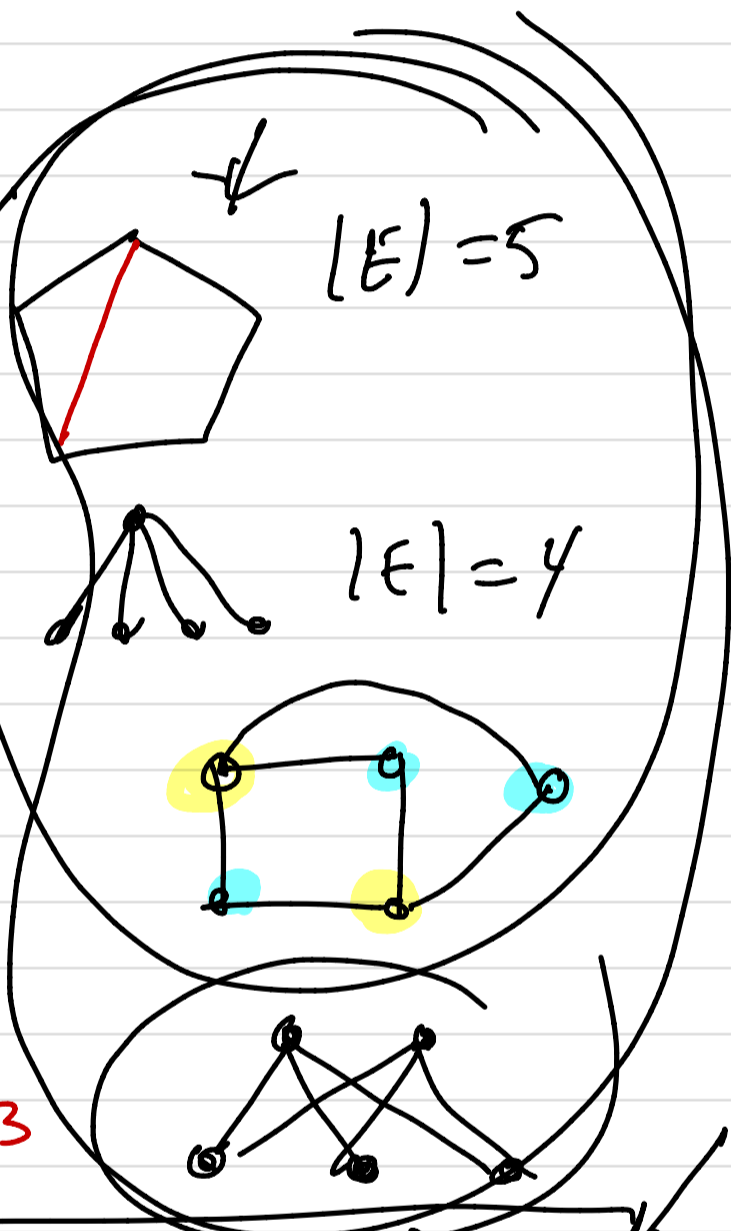
Ex) ?



edges
 4
 3

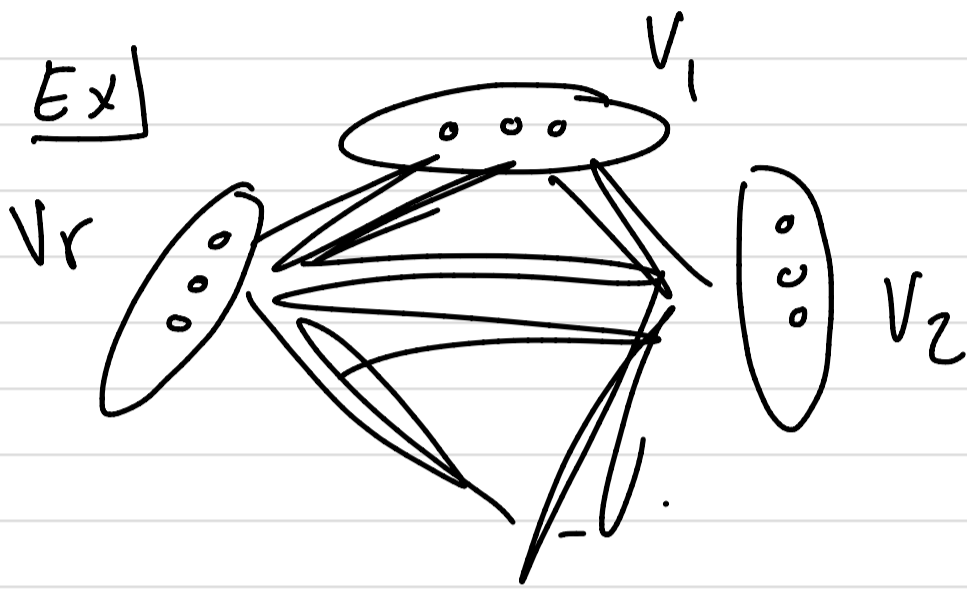


$n=4$
 K^3 critical



$ex(5, K^3) = 6$

How many edges in graph on n vertices that fails to contain a K^r ? Suppose $n \geq r$

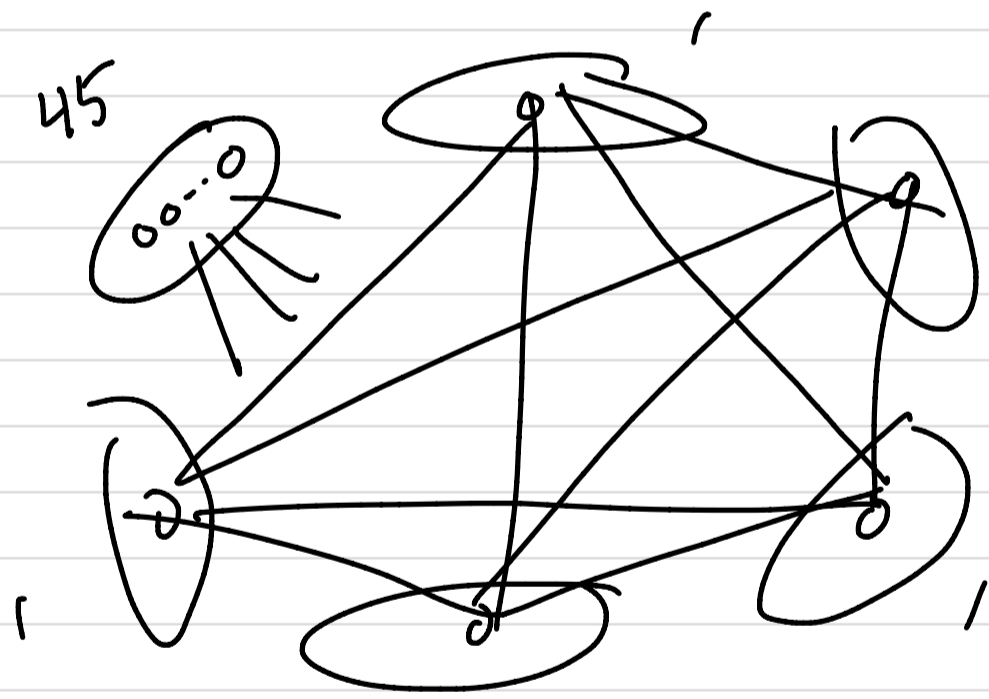


G is $(r-1)$ -partite graph.

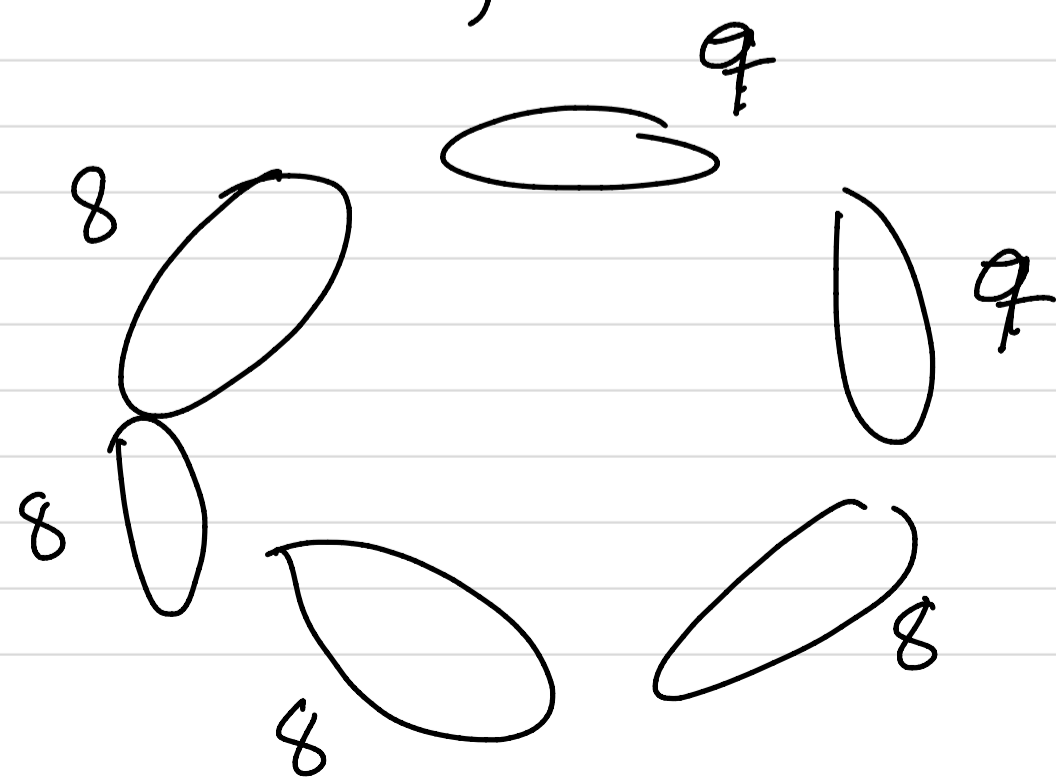
G has no K^r subgraph.

$n=50$ $r=7$ avoiding K^7

Construct a 6-partite graph



$$E(G) = 45 \cdot 5 + \binom{5}{2}$$



$$\begin{array}{r} 32 \\ \del{47} \\ \del{46} \\ \hline 18 \\ \hline 50 \checkmark \end{array}$$

def: $T^r(n)$ an r -partite graph on n vertices such that

the classes are as even as possible.

\forall two class V_i, V_j , $||V_i| - |V_j|| \in \{0, 1\}$.

a Turán graph.

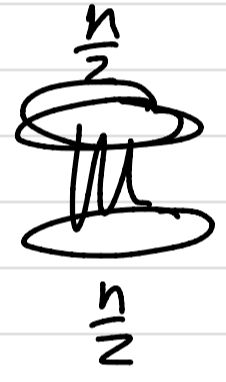
Turán's Thm

$$ex(n, K^{r+1}) = E(T^r(n)) = t_r(n)$$

P_m

$t_2(n) = \begin{cases} \frac{n^2}{4} & \text{even} \\ \frac{n^2-1}{4} & \text{odd} \end{cases}$

n even
 $\sum_1^{n/2} \frac{n}{2}$
 $= \frac{n}{2} \cdot \frac{n}{2}$



$= \lfloor \frac{n^2}{4} \rfloor$

odd

$\frac{(n-1)(n+1)}{4}$
 $= \frac{n^2-1}{4}$



$\lfloor \frac{n^2}{4} \rfloor + 1$ edges